

SINGULAR 3-1-3 – Tackling New Challenges

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Driving force behind SINGULAR's development:
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- cryptographical applications
- commutative and non-comm. Gröbner Bases in group theory

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(\implies e.g. user defined structured data types,
`python_objects`, `python_run`)
- Make use of multicore processors
(\implies parallelization)

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- Regular models for arithmetic surfaces
needs wider functionality over \mathbb{Z} ,
not just Gröbner Bases

Regular models via Lipman's algorithm

Theorem (Lipman)

Let X be an excellent, noetherian, reduced scheme of dimension 2. Then X has a desingularization of the form

$$X_r \xrightarrow{\pi_r \circ n_r} \dots \xrightarrow{\pi_2 \circ n_2} X_1 \xrightarrow{\pi_1 \circ n_1} X_0 = X$$

where each n_i is a normalization, π_i a blow-up at a 0-dim. center.

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apply this over \mathbb{Z} , hence algorithmic tasks:

- normalization
- blow-up (= appropriate Gröbner basis computation)
- primary decomposition (as tool during normalization steps)

Theorem (Grauert-Remmert Criterion for Normality)

R noetherian, reduced ring, $J \subset R$ ideal s.th.

- 1 J contains a non-zerodivisor p on A ,
- 2 J is a radical ideal,
- 3 $V(\mathcal{C}_{\bar{R}|R}) \subset V(J)$.

Then $R = \bar{R} \iff R \cong \text{Hom}_R(J, J) \cong \frac{1}{p}(pJ :_R J) \subset \bar{R} \subset Q(R)$.

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Resulting algorithm requires primary decomposition
(here over \mathbb{Z} to get $\overline{\mathbb{Z}[x]/I}$):

- for equidimensional decomposition,
- for decomposing the singular locus.

Primary Decomposition over \mathbb{Z}

Well-known: primary decomposition in $\mathbb{Q}[\underline{x}]$ and $\mathbb{Z}/p[\underline{x}]$
(e.g. Gianni-Trager-Zacharias)

New aspects over \mathbb{Z} :

- $I \cap \mathbb{Z} = \langle m \rangle$, $m \in \mathbb{Z} \setminus \{0\}$, may occur
for original ideal or at intermediate step

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 - primary components over \mathbb{Q} (in absence of such m)
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- second step above is parallel in nature

On two different levels:

1 SINGULAR-SINGULAR communication

- for larger subtasks (communication overhead)
- improved and extended communication interface
- examples: modStd, primdecZ, modular primary decomposition

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2 Threads

- for fine grained subtasks
- work in progress, long term goal, affects memory management
- examples: polynomial arithmetic, polynomial gcd

Application: Modular Standard Bases

Modular computation of standard basis G of an ideal $I \subset \mathbb{Q}[x]$:

- 1 choose several primes and compute SB G_p modulo each p
- 2 delete unlucky primes
(i.e. $L(I_p) \neq L(I_q)$ for most primes $q \neq p$)
- 3 find G by Chinese remaindering and Farey rational map
- 4 test Gröbner Basis property (most expensive part)
- 5 reiterate with further (new) primes if test fails

Modular Standard Bases: Gröbner Basis test

SINGULAR's method for final verification:

Theorem (Arnold (homogeneous)/ Idrees, Pfister, Steidel (general for global & local orderings))

Let $I \subseteq \mathbb{Q}[\underline{x}]$ an ideal and $G \subseteq I$ a set of polynomials such that

- $L(G) = L(G_p)$ where G_p is a standard basis of I_p for some prime number p ,
- G is a standard basis of $\langle G \rangle$,
- $I \subseteq \langle G \rangle$.

Then $I = \langle G \rangle$.

Modular Standard Bases in Parallel

Parallel steps in modStd:

- 1 Compute G_p for different p in parallel
via SINGULAR-SINGULAR communication
- 2 Parallel final verification tests:
 - test $f \in \langle G \rangle$ for each generator f of I
 - reduce s-polynomials of G w.r.t. Gwill be implemented via threads

Modular Standard Bases and Parallel Computing

Example	std	modStd	modStd ₄ *	modStd ₉ *
cyclic8	-	8271	4120	2927
Paris.ilias13	37734	1159	676	580
homog.cyclic7	3343	3436	886	408
	-	6	3	3

Table: Total running times (in sec) for computing a standard basis of examples chosen from The SymbolicData Project (H.-G. Gräbe) via std, modStd and its parallelized variant modStd_n* for $n = 4, 9$.

Modular Primary Decomposition in Dimension Zero

“Decompose” singular locus of randomly chosen rational plane curves:

degree	minassGTZ	assPrimes
7	39	166
8	193	196
9	3776	344
10	24865	969
11	270886	2535
12	943654	6369

Table: Total running times (in ms), where assPrimes uses 8 cores.

Normalization via Localization

Proposition

R noetherian domain, $\text{Sing}(R) = \{P_1, \dots, P_s\}$, $S_i = R \setminus P_i$ for $i = 1, \dots, s$. Suppose intermediate rings $R \subset R^{(i)} \subset \overline{R}$ are given such that $S_i^{-1}R^{(i)} = \overline{S_i^{-1}R}$. Then $\sum_{i=1}^s R^{(i)} = \overline{R}$.

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Proposition (Local Normality Criterion)

R noetherian domain, $R \subset R'$ module-finite extension, $P \in \text{Sing}(R)$ minimal with respect to inclusion, $S = R \setminus P$, $J' = \sqrt{PR'}$. If $R' \cong \text{Hom}_{R'}(J', J')$, then $S^{-1}R'$ is normal.

Integral Bases, Adjoint Ideals, Parametrization of Rational Curves

R coordinate ring of irreducible plane curve. Then:

- $\text{Sing}(R)$ zero-dimensional \Rightarrow can use modular primary decomposition
- local contributions to normalization also via Puiseux expansions (van Hoeij) and Hensel-lifting (Böhm-Decker-Laplagne-Seelisch)

Important applications: Adjoint ideals, parametrization of rational curves. New, fast algorithms by Böhm-Decker-Laplagne-Seelisch. Two different approaches to compute adjoint ideals via localization.

SINGULAR Developers



Recent Developments since SINGULAR 3-1-2

New contributed/experimental libraries include:

- `integralbasis.lib` – Integral Bases
- `paraplanecurve.lib` – Rational Parametrization of Rational Plane Curves
- `monomialideal.lib` – Operations for Monomial Ideals
- `multigrading.lib` – Multigradings and Related Operations
- `primdecint.lib` – Primary Decomposition over the Integers
- `resbinomial.lib` – Desingularization of Binomial Ideals
- and many more

New/improved features include:

- user defined types
- python objects
- new Singular-Singular communication interfaces